Locality Optimizations in Tapestry

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Why Is This a Problem?

- Example Application: OceanStore web caching
 - If a nearby replica exists, we **must** find it quickly
- Measure of Locality: Relative Delay Penalty (RDP)
 - The ratio of the distance of an object in Tapestry to the minimum possible distance (i.e. over IP)
- **Problem**: finding nearby objects incurs a high RDP
 - Two extra hops have a huge relative impact if object is close
 - An issue for all similar systems, not just Tapestry
- Solution: trade storage overhead for low RDP

Optimization 1: Publish to Backups

• **Redundancy:** Routing table entries store up to *c* nodes

- Closest node is the *primary neighbor*, *c*-1 nodes are *backups*

• A simple optimization: publish to *k* backups

– Limit to the first *n* hops of the publish path

Result

 Nodes near the object more likely to encounter pointers while routing to the root

- Storage overhead: k*n additional pointers per object



Optimization 1: Publish to Backups

n = 2 hops





Optimization 2: Local Misroute

- Solution: Before taking "long" hop, misroute to closer node
 - Look a little harder in the local area before leaving
 - When publishing, place a pointer on *local surrogate*
- Issue: What determines a "long" hop?
 - One metric: if next hop is more than *m* times longer than last hop, consider it "long"
 - Call *m* the threshold factor

Optimization 2: Local Misroute



Experiments run in simulation on a transit-stub topology "Using sim knowledge" indicates direct use of the topology file

RDP

Future Work

- Analyze more locality optimizations and different
 parameter configurations
- Take measurements on PlanetLab
- Test optimizations with real workloads (i.e. web caching)
- Complete cost analysis of storage overhead vs. RDP benefit across all optimizations