#### Optimizations for Locality-Aware Structured Peer-to-Peer Overlays

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## Discussion

#### • Why is this a problem?

- Latency, efficiency, availability

Metric: Relative Delay Penalty (RDP)
 – Distance through Tapestry vs. IP distance

Solution: trade storage for low local area RDP
 Will work in DOLRs with a pointer indirection layer

# **Optimization 1: Backups**

- Redundancy: Store up to c nodes in each entry
  - c-1 nodes are backups
- A simple optimization: publish to b backups
  - Limit to first h hops of publish path
- Result

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- Nodes near the object more likely to encounter pointers
- Cost: b\*h additional pointers per object



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### **Optimization 1: Backups**



## **Optimization 2: Nearest Neighbors**

- Observation: In Opt. 1, choice for backups is limited
  - But lots of nodes at each level, many may be nearby
- Optimization: publish to n nearest neighbors
  - Limit to first h hops of the publish path
- Result
  - If n is large, essentially local area flooding
  - Analytical cost: *n*\**h* additional pointers per object



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### **Optimization 2: Nearest Neighbors**





### **Optimization 3: Local Surrogate**

- Solution: Check local node before leaving
  - When publishing, place a pointer on *local surrogate*
  - Occurs naturally on Coral, LAND, SkipNet
  - In practice, storage cost is very low
- Issue: What determines a wide area hop?
  One metric: if next hop is more than t times longer than last hop, consider it wide area



#### **Optimization 3: Local Surrogate**



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## **Future Work**

- Automatically adjust t when using local surrogate
- Take measurements on actual networks
- Test optimizations with real workloads
- Evaluate the maintenance cost

#### **Questions?**