

6.5081 : VM for user applications

OS kernel use VM in creative ways

Paper argues user apps can use VM too

- Garbage collector
- Data compression
- SVM

What primitives?

Trap → alarm handler → sigaction

protect Prot1 — decrease accessibility
R+W → R → rd access

ProtN — some TLB flushes

Unprot — increase accessibility

? Dirty ? Map2 — several mmap's

Unix today. mmap

map a file

addr = mmap(NULL, len, RW, MAP_PRIVATE,
fd, offset)

Next lab

map anonymous

Unix today

mprotect (addr, len, R) → lds
NONE

→ sigalarm

unmap: remove address range

sigaction: signal handler (f)
signal: segfault.

VM implementation

AS : page table + Virtual Memory Area (VMAs)

→ contiguous range of addresses
same permission
backed by same object

1000-2000

2100-2200

User-level traps

- PTE marked invalid / R \emptyset
- CPU jumps ^{in kernel}
- Kernel save state
- ASKs the VM system what to do? VMMS
- Upcall into user space
- Run handler \rightarrow improtect?
- handler returns to kernel.
- Kernel resumes interrupted process

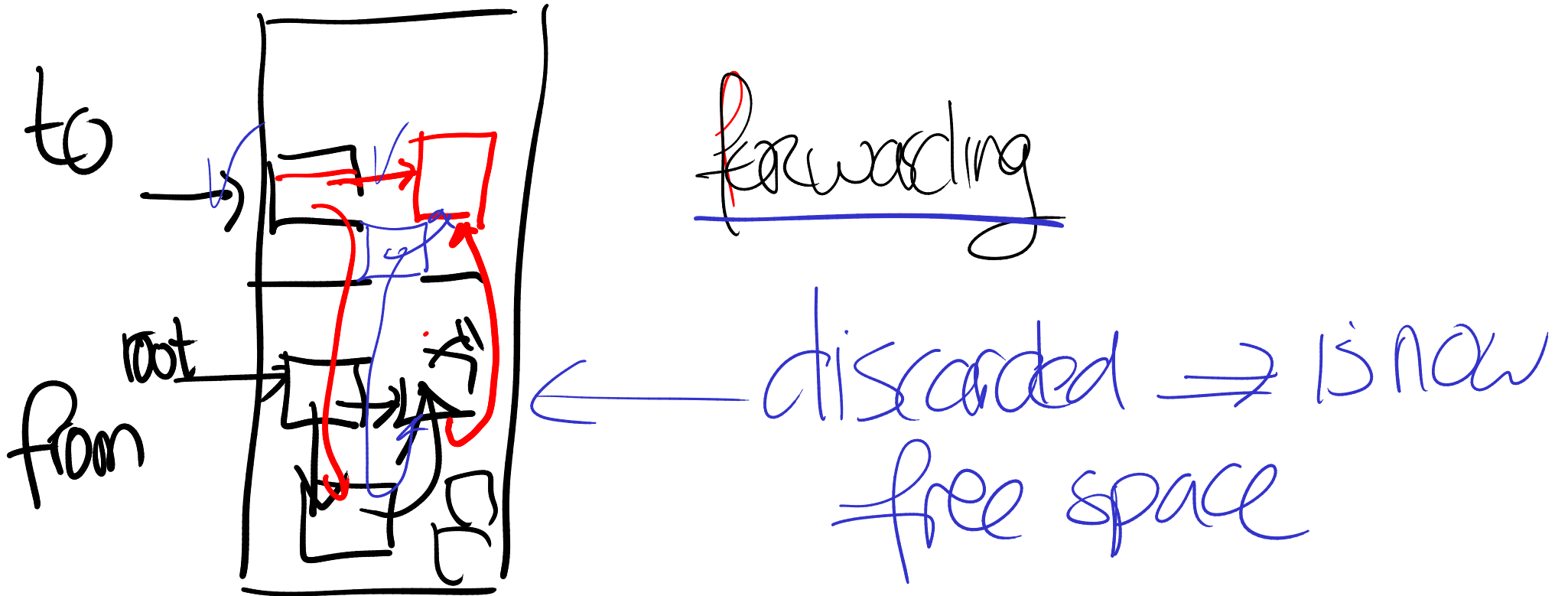
Example: Huge memorization table



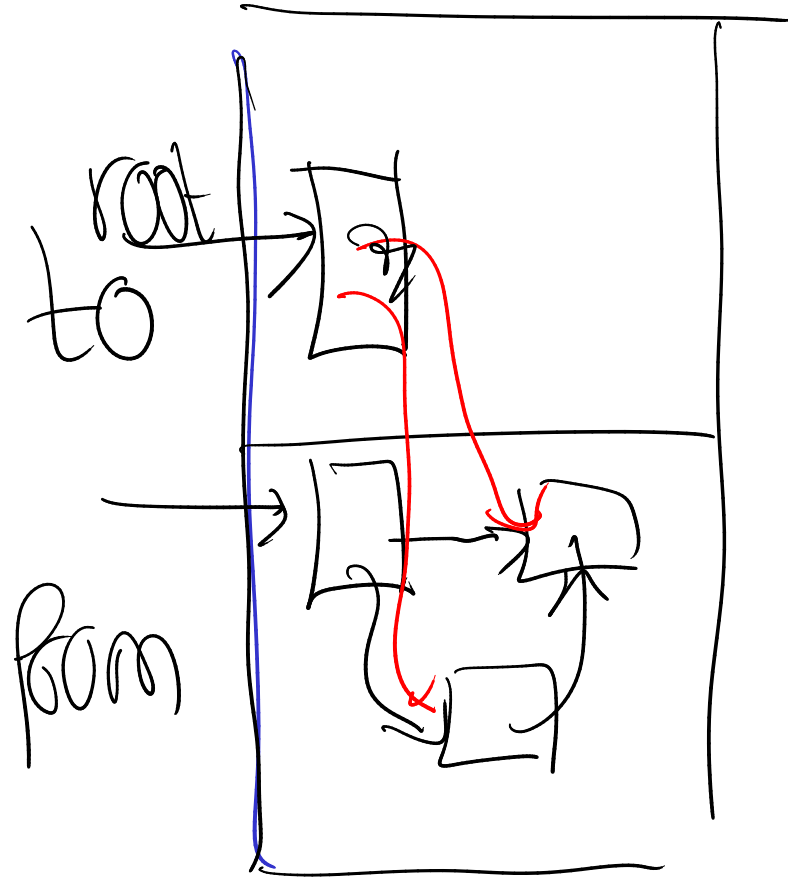
$f(i)$? \rightarrow lookup
[h[i]]

Example: garbage collector

A copying garbage collector



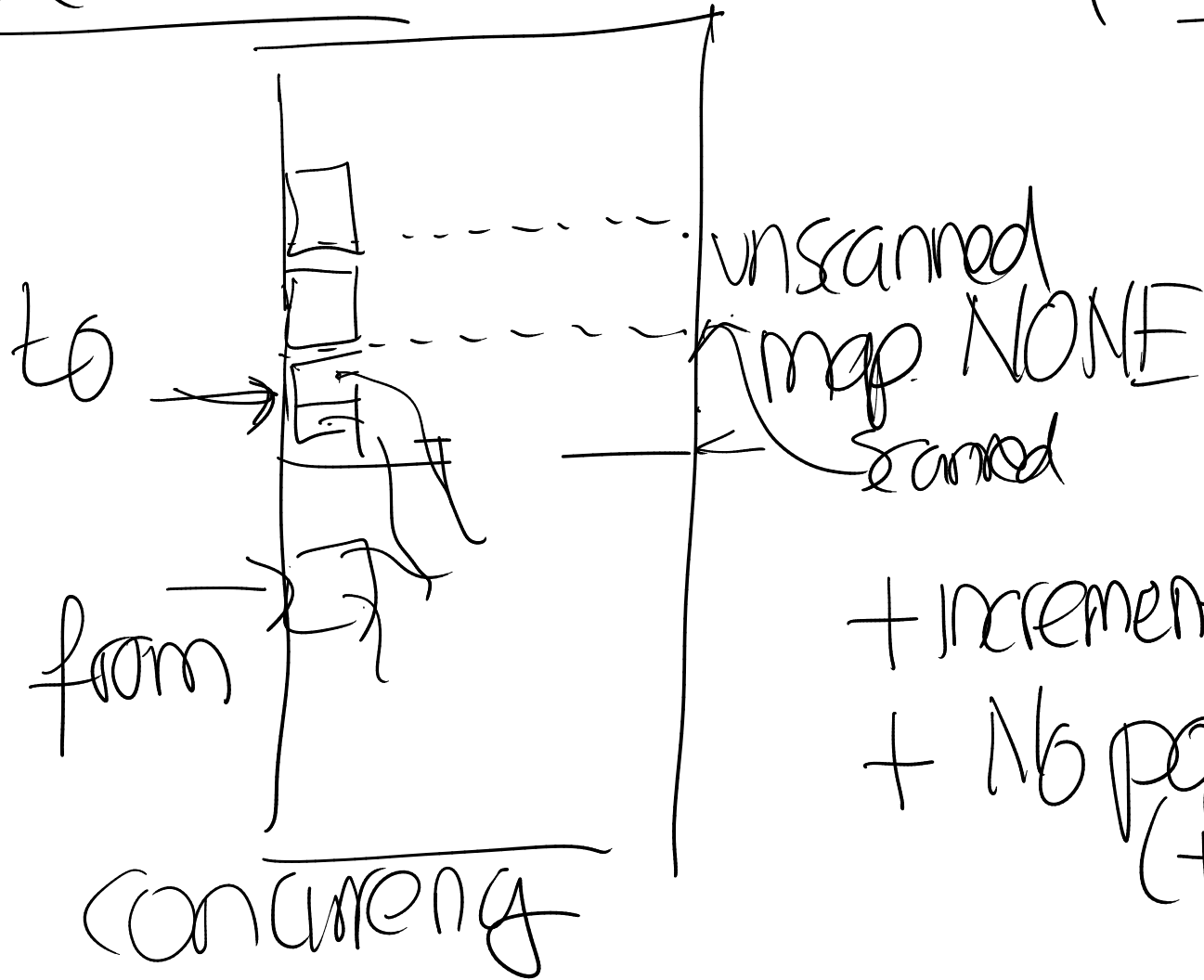
Baker's: real-time & incremental gc



new: forward a few more objects

dereference a pointer
check if in from space
=> forwarding

Use VM:



fault handler

Scan one page
of objects +
forward them

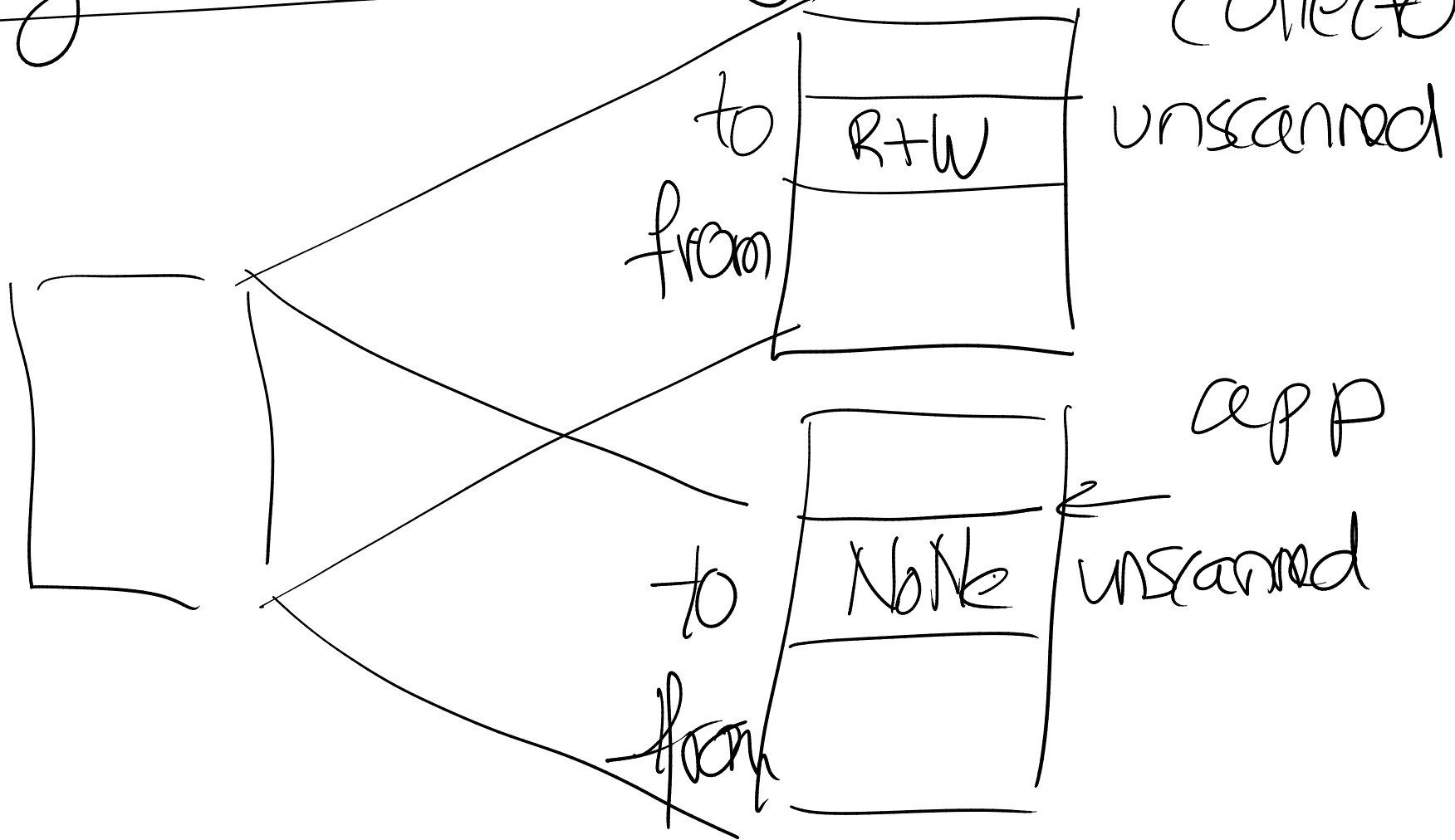
unprotect
the scanned page

+ incremental

+ No pointer check.

(the VM hw
does it for us)

Risky issue w. concurrency: maps collectors



Should use VM?

Most cases could have been
implemented w.
extra install

↳ ~~computer~~

Checkpointing
Sum

Unix
supports them.

What has changed 1991?

Many!

Continuous development

Some big changes:

5-levels

ASIO

KPTI